

Fig. 3. Graphical representation of conditions (33) and (34). R.P. is guaranteed since $|S|<c_{s}$ for $\omega<2$ and $|H|<c_{H}$ for $\omega>1.4$.

The step from (A1) to (A2) follows from $M=N_{11}+N_{12}(I-$ $\left.T N_{22}\right)^{-1} T N_{21}$ and Schurs formula

$$
\operatorname{det}\left(A-B D^{-1} C\right)=\operatorname{det}\left[\begin{array}{cc}
A & B  \tag{A3}\\
C & D
\end{array}\right] / \operatorname{det} D
$$

and the assumption $\operatorname{det} D=\operatorname{det}\left(I-T N_{22}\right) \neq 0$.
Lemma 1: An equivalent statement of the lemma is as follows: Let $\tilde{\Delta}$ $=\operatorname{diag}\left\{\Delta_{1}, \Delta_{2}\right\}$ where $\Delta_{1}$ and $\Delta_{2}$ have the same size as $N_{11}$ and $N_{22}$, respectively. ( $\Delta_{1}$ and $\Delta_{2}$ may have additional structure.) Then

$$
\mu_{\bar{\Lambda}}\left[\begin{array}{cc}
0 & N_{12}  \tag{A4}\\
c N_{21} & 0
\end{array}\right]=\sqrt{c} \mu_{\bar{\Omega}}\left[\begin{array}{cc}
0 & N_{12} \\
N_{21} & 0
\end{array}\right] .
$$

Proof of (A4):

$$
\begin{gather*}
\mu_{\bar{\Delta}}\left[\begin{array}{cc}
0 & N_{12} \\
c N_{21} & 0
\end{array}\right] \leq 1 / k_{1}  \tag{A5}\\
\Leftrightarrow \operatorname{det}\left(I+k_{1}\left[\begin{array}{cc}
\Delta_{1} & \\
& \Delta_{2}
\end{array}\right]\left[\begin{array}{cc}
0 & N_{12} \\
c N_{21} & 0
\end{array}\right]\right) \neq 0 \\
\Leftrightarrow \operatorname{det}\left[\begin{array}{cc}
I & k_{1} \Delta_{1} N_{12} \\
k_{1} c \Delta_{2} N_{21} & I
\end{array}\right] \neq 0  \tag{A6}\\
\Leftrightarrow \operatorname{det}\left(I-k_{1}^{2} c \Delta_{1} N_{12} \Delta_{2} N_{21}\right) \neq 0 \\
\Leftrightarrow \operatorname{det}\left[\begin{array}{cc}
I & \sqrt{k_{1}^{2} c} \Delta_{1} N_{12} \\
\sqrt{k_{1}^{2} c} \Delta_{2} N_{21} & I
\end{array}\right] \neq 0  \tag{A8}\\
\Leftrightarrow \mu_{\bar{\Delta}}\left[\begin{array}{cc}
0 & N_{12} \\
N_{21} & 0
\end{array}\right] \leq 1 / \sqrt{k_{1}^{2} c} \tag{A9}
\end{gather*}
$$

The conditions involving det () $\neq 0$ must hold for $\forall \Delta_{1}$ s.t. $\bar{\sigma}\left(\Delta_{1}\right)<1$ and $\forall \Delta_{2}$ s.t. $\bar{\sigma}\left(\Delta_{2}<1\right)$. The step from (A6) to (A7) and back to (A8) follows from (A3). Since (A5) and (A9) must hold for any value of $k_{1}$, (A4) follows.
Theorem 2: From Theorem 1 and Lemma 1 for the case $N_{11}=N_{22}=$ 0

$$
\mu_{\Delta}\left(N_{12} T N_{21}\right) \leq k \text { if } \bar{\sigma}(T) \mu_{\Sigma}^{2}\left[\begin{array}{cc}
0 & N_{12}  \tag{A10}\\
N_{21} & 0
\end{array}\right] \leq k
$$

Since (A10) must hold for any choice of $k$ it is equivalent to

$$
\mu_{\Delta}\left(N_{12} T N_{21}\right) \leq \tilde{\sigma}(T) \mu_{\tilde{\Delta}}^{2}\left[\begin{array}{cc}
0 & N_{12} \\
N_{21} & 0
\end{array}\right] .
$$

Theorem 2 follows by choosing $N_{12}=A, N_{21}=B$.
Special Cases of Theorem 2: Let $\Delta_{1}$ and $\Delta_{2}$ have the same structure as $\Delta$ and $T$ in Theorem 2. Define $\tilde{\Delta}=\operatorname{diag}\left\{\Delta_{1}, \Delta_{2}\right\}$. Then

$$
\begin{align*}
& \mu_{\tilde{\Delta}}^{2}\left[\begin{array}{cc}
0 & A \\
B & 0
\end{array}\right] \leq 1 / k  \tag{A11}\\
& \Leftrightarrow \mu_{\tilde{\Lambda}}\left[\begin{array}{cc}
0 & k A \\
B & 0
\end{array}\right] \leq 1 \tag{A12}
\end{align*}
$$

$$
\begin{gather*}
\Leftrightarrow \operatorname{det}\left(I+\left[\begin{array}{ll}
\Delta_{1} & \\
& \Delta_{2}
\end{array}\right]\left[\begin{array}{cc}
0 & k A \\
B & 0
\end{array}\right]\right) \neq 0 \forall \Delta_{1}, \Delta_{2} \\
\Leftrightarrow \operatorname{det}\left(I-k \Delta_{2} B \Delta_{1} A\right)=\operatorname{det}\left(I-k \Delta_{1} A \Delta_{2} B\right) \neq 0 \quad \forall \Delta_{1}, \Delta_{2} \\
\Leftrightarrow \mu_{\Delta_{2}}\left(B \Delta_{1} A\right) \leq 1 / k \forall \Delta_{1}  \tag{A13}\\
\Leftrightarrow \mu_{\Delta_{1}}\left(A \Delta_{2} B\right) \leq 1 / k \forall \Delta_{2}  \tag{A14}\\
\Leftrightarrow \rho\left(\Delta_{1} A \Delta_{2} B\right)=\rho\left(\Delta_{2} B \Delta_{1} A\right) \leq 1 / k \forall \Delta_{1}, \Delta_{2} . \tag{A15}
\end{gather*}
$$

By $\forall \Delta_{i}$ is understood all $\Delta_{i}$ s.t. $\bar{\sigma}\left(\Delta_{i}\right)<1$. The step from (A11) to (A12) follows from (A4).

Case 1): Follows from (A15): Use the SVD of $A=U_{A} \Sigma_{A} V_{A}^{H}$ and $B$ $=U_{B} \Sigma_{B} V_{B^{-}}^{H}$. Since $\Delta_{1}$ and $\Delta_{2}$ are full, $\Delta_{1}$ may be chosen such that $\Delta_{1} U_{A}$ $=V_{B}$ and $\Delta_{2}$ such that $V_{A}^{H} \Delta_{2}=U_{B}^{H}$. Then $\rho\left(\Delta_{1} A \Delta_{2} B\right)=\rho\left(V_{B} \Sigma_{1} \Sigma_{2} V_{B}^{H}\right)$ $=\rho\left(\Sigma_{1} \Sigma_{2}\right)=\bar{\sigma}(A) \bar{\sigma}(B)$. [The generalization to the case when $A$ and $B$ are nonsquare is straightforward and involves lining up the directions corresponding to $\bar{\sigma}(A)$ and $\bar{\sigma}(B)$.]

Case 2): Follows from (A14).
Case 3): Follows from (A13).
Cases 4), 5): Follow from (A15).

## AcKNOWLEDGMENT

The authors are thankful to J. C. Doyle for numerous useful discussions and remarks. The idea of treating $T$ as a perturbation (which subsequently led to the derivation of Theorem 1) was first presented by Grosdidier and Morari [4] in their derivation of the $\mu$-interaction measure.

## REfERENCES

[1] J. C. Doyle, "Analysis of feedback systems with structured uncertainties," Proc. IEE, Pt. D., vol. 129, pp. 242-247, 1982.
[2] J. C. Doyle, J. E. Wall, and G. Stein, "Performance and robustness analysis for structured uncertainty," Proc. IEEE Conf. Decision Contr., Orlando, FL, 1982.
[3] J. C. Doyle, Lecture Notes for ONR/Honeywell Workshop, Minneapolis, MN, Oct. 8-10, 1984.
[4] P. Grosdidier and M. Morari, "Interaction measures for systems under decentralized control,"' Automatica, vol. 22, no. 3, pp. 309-319, 1986.
[5] I. Postethwaite and Y. K. Foo, "Robustness with simultaneous pole and zero movements across the j $\omega$-axis," Automatica, vol. 21, no.4, pp. 433-443, 1985.

## Recursive Algorithm for the Computation of the $H^{\infty}$-Norm of Polynomials

## LEI GUO, LIGE XIA, AND YI LIU

Abstract-A recursive algorithm for computing the $H^{\infty}$-norm of polynomials is developed. The algorithm is shown to converge monotonically and the convergence rate is also established. Some examples are presented to illustrate the algorithm.

## I. INTRODUCTION

In recent years, $H^{\infty}$-norm and its optimization have been used more and more frequently in many areas of control theory and applications. For example, $H^{\infty}$-norm optimal controller synthesis approach [1], [2], model/

[^0]controller reduction, and even some problems in system identification are closely related to $H^{\infty}$-norms. The model/controller reduction is often best posed as a frequency weighted $H^{\infty}$ optimal approximation problem [3]. For a given transfer function $G(z)$, many approaches give a reduced-order transfer function $G_{r}(z)$, normally, which is not optimal in the $H^{\infty}$ sense. Certainly, it is desirable to know the value of the $H^{\infty}$-approximation error $\left\|G(z)-G_{r}(z)\right\|_{\infty}$. In system identification, if a monic polynomial $C(z)$ is the moving average noise process transfer function in an ARMAX model, it is well known that for the convergence of the extended least-squares algorithm, a key condition is that $C^{-1}(z)-1 / 2$ is strictly positive real (e.g., [4], [5]). It is easy to see that this condition is equivalent to the requirement $\|C(z)-1\|_{\infty}<1$. However, in practice, to calculate the value of the $H^{\infty}$-norm is not a pleasant task. It is usually done by a rather trivial method, i.e., plotting the absolute value of the function concerned on the unit circle.

In this note, we propose a theoretical recursive algorithm for the computation of the $H^{\infty}$-norm of polynomials or FIR transfer functions (Section II). In Section III we give the derivation of the algorithm and show that the guaranteed convergence rate of the algorithm is $O(\log n / n)$. Simulation results of some examples are provided in Section IV. Section V concludes the note with some remarks.

Before pursuing further, we need some concepts and definitions as follows.

Let $f(z)$ be a complex-valued function on the unit circle bounded almost everywhere; the set of all such functions is denoted by $L^{\infty}$, with norm

$$
\begin{equation*}
\|f(z)\|_{\infty}=\underset{\theta \in[0,2 \pi]}{\operatorname{ess} \sup }\left|f\left(e^{i \theta}\right)\right| . \tag{1}
\end{equation*}
$$

The Hardy space $H^{\infty}$ consists of all complex-valued functions which are analytic and of bounded modulus on $|z|<1$, with norm

$$
\begin{equation*}
\|f(z)\|_{\infty}=\sup _{|z|<1}|f(z)| \tag{2}
\end{equation*}
$$

It is known that each $f$ in $H^{\infty}$ yields a unique $L^{\infty}$ boundary function with the two norms equal. The set of such boundary functions is the subspace of $L^{\infty}$-functions with Fourier coefficients zero for negative indexes, and we can regard $H^{\infty}$ as a closed-subspace of space $L^{\infty}$.
We also need the concept of space $L^{p}(p>0)$. It consists of all measurable complex functions $f(z)$ defined on the unit circle $|z|=1$ such that $\left|f\left(e^{i \theta}\right)\right| p$ is integrable with respect to Lebesgue measure, with norm

$$
\begin{equation*}
\|f(z)\|_{p}=\left\{\frac{1}{2 \pi} \int_{0}^{2 \pi}\left|f\left(e^{i \theta}\right)\right|^{p} d \theta\right\}^{1 / p} . \tag{3}
\end{equation*}
$$

## II. Algorithm Description and Main Results

Let $C(z)$ be a polynomial with real coefficients and with degree $r$

$$
\begin{equation*}
C(z)=C_{0}+C_{1} z+\cdots+C_{r} z^{r}, \quad C_{0} C_{r} \neq 0 \tag{4}
\end{equation*}
$$

Define a function $f(z)$ as

$$
\begin{align*}
f(z) & =C(z) C\left(z^{-1}\right) \\
& \triangleq \gamma_{0}+\sum_{j=1}^{r} \gamma_{j}\left(z^{j}+z^{-j}\right) \tag{5}
\end{align*}
$$

where

$$
\begin{equation*}
\gamma_{j}=\sum_{k=0}^{r} C_{k} C_{k+j}, \quad\left(C_{k}=0, \text { for } k>r\right) \tag{6}
\end{equation*}
$$

To describe our algorithm, we need the following auxiliary variables:

$$
\left\{\chi_{i}(n), 1 \leqslant i \leqslant 2 r, n \geqslant 1\right\} \text { and }\{T(n), n \geqslant 1\}
$$

which are recursively defined by

$$
\begin{align*}
& \chi_{k+r}(n-1)= {\left[\sum_{j=1}^{r}(n j-k) \gamma_{j} \chi_{k-j}(n-1)\right.} \\
&\left.-\sum_{j=0}^{r-1}(n j+k) \gamma_{j} \chi_{k+j}(n-1)\right] /(n r+k) \gamma_{r} \quad 1 \leqslant k \leqslant r  \tag{7}\\
& \chi_{k}(n)=(n / k) \sum_{j=1}^{r} j \gamma_{j}\left[\chi_{k-j}(n-1)-\chi_{k+j}(n-1)\right] \\
& \cdot {\left[\gamma_{0}+2 \sum_{j=1}^{r} \gamma_{j} \chi_{j}(n-1)\right]{ }^{-1} \quad 1 \leqslant k \leqslant r }  \tag{8}\\
& T(n)=\frac{n-1}{n} T(n-1)+\frac{1}{2 n} \log \left[\gamma_{0}+2 \sum_{j=1}^{r} \gamma_{j} \chi_{j}(n-1)\right] \tag{9}
\end{align*}
$$

where by definition

$$
\chi_{0}(n)=1 \text { and } \chi_{-i}(n)=\chi_{i}(n), \quad 1 \leqslant i \leqslant 2 r, n \geqslant 1
$$

and where the initial conditions are

$$
\chi_{j}(1)=\gamma_{j} / \gamma_{0}, 1 \leqslant j \leqslant r ; \quad T(1)=\frac{1}{2} \log \gamma_{0} .
$$

The $n$th approximation for the norm $\|C(z)\|_{\infty}$ is defined by

$$
\begin{equation*}
J(n)=\exp \{T(n)\}, \quad n \geqslant 1 . \tag{11}
\end{equation*}
$$

The asymptotic properties of the above algorithm are summarized in the following theorem.

Theorem 1: For any polynomial $C(z)$ defined as in (4), the quantity $J(n)$ given by (7)-(11) increases monotonically and converges to $\|C(z)\|_{\infty}$ as $n \rightarrow \infty$, with convergence rate

$$
\begin{equation*}
\|C(z)\|_{\infty}-J(n) \leqslant\left(\|C(z)\|_{\infty}\right) \frac{\log n}{2 n}+0(1 / n) \tag{12}
\end{equation*}
$$

## III. CONVERGENCE ANALYSIS

For the proof of Theorem 1, we first establish the following lemmas. Lemma I: For $T(n)$ given by (9)

$$
T(n)=\log \left(\|C(z)\|_{2 n}\right)
$$

holds for any $n \geqslant 1$.
Proof: Define

$$
\begin{equation*}
M_{k}(n)=\frac{1}{2 \pi} \int_{0}^{2 \pi} f^{n}\left(e^{i \theta}\right) e^{k i \theta} d \theta \tag{13}
\end{equation*}
$$

for $n \geqslant 1,-2 r \leqslant k \leqslant 2 r$, where $f\left(e^{i \theta}\right)$ is given by (5).
It is easy to see that for any $n \geqslant 1$

$$
\begin{equation*}
M_{-k}(n)=M_{k}(n), \quad k=0,1, \cdots, 2 r \tag{14}
\end{equation*}
$$

and

$$
\begin{align*}
M_{0}(n) & =\frac{1}{2 \pi} \int_{0}^{2 \pi} f^{n}\left(e^{i \theta}\right) d \theta \\
& =\left\|f^{n}(z)\right\|_{n}^{n}=\|C(z)\|_{2 n}^{2 n} \tag{15}
\end{align*}
$$

So for the proof of the lemma, we need only to show that

$$
\begin{equation*}
T(n)=\frac{1}{2 n} \log M_{0}(n) \tag{16}
\end{equation*}
$$

We proceed as follows.

By (5), (13), and (14)

$$
\begin{align*}
M_{0}(n) & =\frac{1}{2 \pi} \int_{0}^{2 \pi} f^{n-1}\left(e^{i \theta}\right) f\left(e^{i \theta}\right) d \theta \\
& =\frac{1}{2 \pi} \int_{0}^{2 \pi} f^{n-1}\left(e^{i \theta}\right)\left[\gamma_{0}+\sum_{k=1}^{r} \gamma_{k}\left(e^{k i \theta}+e^{-k i \theta}\right)\right] d \theta \\
& =\gamma_{0} M_{0}(n-1)+2 \sum_{j=1}^{r} \gamma_{j} M_{j}(n-1) \\
& =M_{0}(n-1)\left\{\gamma_{0}+2 \sum_{j=1}^{r} \gamma_{j}\left[M_{j}(n-1) / M_{0}(n-1)\right]\right\} \tag{17}
\end{align*}
$$

consequently, we have

$$
\begin{aligned}
{\left[\frac{1}{2 n} \log M_{0}(n)\right]=\frac{n-1}{n} } & {\left[\frac{1}{2(n-1)} \log M_{0}(n-1)\right] } \\
& +\frac{1}{2 n} \log \left\{\gamma_{0}+2 \sum_{j=1}^{r} \gamma_{j}\left[M_{i}(n-1) / M_{0}(n-1)\right]\right\}
\end{aligned}
$$

Comparing this to (9), we see that for (16) it suffices to show that

$$
\begin{equation*}
\chi_{j}(n)=M_{j}(n) / M_{0}(n), \quad 1 \leqslant j \leqslant r . \tag{18}
\end{equation*}
$$

Now, by integral parts from (13) and the fact that $f(z)=f\left(z^{-1}\right)$, we have

$$
\begin{align*}
M_{k}(n)= & \frac{1}{2 \pi} \int_{0}^{2 \pi} f^{n}\left(e^{-i \theta}\right) e^{k i \theta} d \theta \\
= & \frac{1}{2 \pi k i} \int_{0}^{2 \pi} f^{n}\left(e^{-i \theta}\right) d e^{k i \theta} \\
= & \frac{1}{2 \pi k i}\left\{\left.f^{n}\left(e^{-i \theta}\right) e^{k i \theta}\right|_{\theta} ^{2 \pi}\right. \\
& \left.-\int_{0}^{2 \pi}\left(e^{k i \theta}\right) d\left[f^{n}\left(e^{-i \theta}\right)\right]\right\} \\
= & \frac{n}{2 \pi k} \int_{0}^{2 \pi} f^{n-1}\left(e^{-i \theta}\right) \cdot f^{\prime}\left(e^{-i \theta}\right) e^{(k-1) i \theta} d \theta \tag{19}
\end{align*}
$$

where

$$
\begin{align*}
f^{\prime}\left(e^{-i \theta}\right) & \left.\frac{d f(z)}{d z}\right|_{z=e^{-i \theta}} \\
& =\sum_{j=1}^{\prime} j \gamma_{j}\left[e^{(1-j) i \theta}-e^{(j+1) i \theta}\right] \tag{20}
\end{align*}
$$

For (19) and (20) we immediately have ( $1 \leqslant k \leqslant r$ )

$$
\begin{equation*}
M_{k}(n)=\frac{n}{k} \sum_{j=1}^{\prime} j \gamma_{j}\left[M_{k-j}(n-1)-M_{k+j}(n-1)\right] \tag{21}
\end{equation*}
$$

Multiplying $1 / M_{0}(n)$ on both sides of this equality and using (17), we know that the recursion (8) is true with $\chi_{k}(n)$ replaced by $M_{k}(n) / M_{0}(n)$.
To conclude (18), we still need to show that the recursion (7) also holds with $\chi_{k}(n)$ replaced by $M_{k}(n) / M_{0}(n)$. To this end, consider the following decomposition for $f^{\prime}\left(e^{-i \theta}\right)$ :

$$
\begin{equation*}
f^{\prime}\left(e^{-i \theta}\right)=g_{1}\left(e^{-i \theta}\right)-r f\left(e^{-i \theta}\right) e^{i \theta} \tag{22}
\end{equation*}
$$

where

$$
\begin{equation*}
g_{1}\left(e^{-i \theta}\right)=r \gamma_{0} e^{i \theta}+\sum_{j=1}^{r} \gamma_{j}\left[(r+j) e^{(1-j) i \theta}+(r-j) e^{(j+1) i \theta}\right] \tag{23}
\end{equation*}
$$

Substituting (22) into (19), we get

$$
\begin{aligned}
M_{k}(n)= & \frac{n}{2 \pi k} \int_{0}^{2 \pi} f^{n-1}\left(e^{-i \theta}\right)\left[g_{1}\left(e^{-i \theta}\right)\right. \\
& \left.-r f\left(e^{-i \theta}\right) e^{i \theta}\right] e^{(k-1) i \theta} d \theta \\
= & \frac{-n r}{2 \pi k} \int_{0}^{2 \pi} f^{n}\left(e^{-i \theta}\right) e^{k i \theta} d \theta \\
+ & \frac{n}{2 \pi k} \int_{0}^{2 \pi} f^{n-1}\left(e^{-i \theta}\right) g_{1}\left(e^{-i \theta}\right) e^{(k-1) i \theta} d \theta \\
= & \frac{-n r}{k} M_{k}(n)+\frac{n}{2 \pi k} \int_{0}^{2 \pi} f^{n-1}\left(e^{-i \theta}\right) \\
& \cdot g_{1}\left(e^{-i \theta}\right) e^{(k-1) i \theta} d \theta .
\end{aligned}
$$

By this identity, we obtain for $1 \leqslant k \leqslant r$

$$
\begin{align*}
M_{k}(n)= & \frac{n}{n r+k} \cdot \frac{1}{2 \pi} \int_{0}^{2 \pi} f^{n-1}\left(e^{-i \theta}\right) g_{1}\left(e^{-i \theta}\right) e^{(k-1) / \theta} d \theta \\
= & \frac{n}{n r+k}\left\{r \gamma_{0} M_{k}(n-1)+\sum_{j=1}^{r} \gamma_{j}\left[(r+j) M_{k-j}(n-1)\right.\right. \\
& \left.\left.+(r-j) M_{k+j}(n-1)\right]\right\} \tag{24}
\end{align*}
$$

which in conjunction with (21) gives the recursive formula for $M_{k+r}(n-$ 1)

$$
\begin{aligned}
M_{k+r}(n-1)= & \frac{1}{(n r+k) \gamma_{r}}\left[\sum_{j=i}^{r}(n j-k) \gamma_{j} M_{k-j}(n-1)\right. \\
& \left.\quad-\sum_{j=0}^{r-1}(n j+k) \gamma_{j} M_{k+j}(n-1)\right] .
\end{aligned}
$$

From here it is easy to see that (7) is true with $\chi_{k}(n-1)$ replaced by $M_{k}(n-1) / M_{0}(n-1)$. This proves the assertion (18), and hence the conclusion of the lemma.
Lemma 2: Let a complex function $f(z) \in L^{\infty}$, if $d / d \theta\left[\left|f\left(e^{i \theta}\right)\right|^{2}\right] \epsilon L^{\infty}$; then

$$
0 \leqslant\left\|f\left(e^{i \theta}\right)\right\|_{\infty}-\left\|f\left(e^{i \theta}\right)\right\|_{n} \leqslant\left(\left\|f\left(e^{i \theta}\right)\right\|_{\infty}\right) \frac{\log n}{n}+O\left(\frac{1}{n}\right) .
$$

Proof: By (1) and (3) it is evident that for any $n \geqslant 1$

$$
\left\|f\left(e^{i \theta}\right)\right\|_{n} \leqslant\left\|f\left(e^{i \theta}\right)\right\|_{\infty} .
$$

Now, denote

$$
g(\theta)=\left|f\left(e^{i \theta}\right)\right|^{2}, \quad \theta \in[0,2 \pi]
$$

Since $g(\theta)$ is a continuous function of $\theta$, there exists a $\theta_{0} \in[0,2 \pi]$ such that

$$
g\left(\theta_{0}\right)=\max _{\theta \in[0,2 \times]} g(\theta)=\left\|f\left(e^{i \theta}\right)\right\|_{\infty}^{2}
$$

Without loss of generality, assume $\theta_{0} \epsilon(0,2 \pi)$. By the Taylor expansion we know that

$$
g(\theta)=g\left(\theta_{0}\right)+g^{\prime}(\xi)\left(\theta-\theta_{0}\right)
$$

where $\xi$ is some point between $\theta$ and $\theta_{0}$.

From here we have, for sufficiently large $n$,

$$
\begin{aligned}
&\left\|f\left(e^{i \theta}\right)\right\|_{n} \\
&=\left\{\frac{1}{2 \pi} \int_{0}^{2 \pi}\left|f\left(e^{i \theta}\right)\right|^{n} d \theta\right\}^{1 / n} \\
&=\left\{\frac{1}{2 \pi} \int_{0}^{2 \pi}[g(\theta)]^{n / 2} d \theta\right\}^{1 / n} \\
&=\left\{\frac{1}{2 \pi} \int_{0}^{2 \pi} g^{n / 2}\left(\theta_{0}\right)\left[1+\frac{g^{\prime}(\xi)}{g\left(\theta_{0}\right)}\left(\theta-\theta_{0}\right)\right]^{n / 2} d \theta\right\}^{1 / n} \\
&=\left[g\left(\theta_{0}\right)\right]^{1 / 2}\left\{\frac{1}{2 \pi} \int_{0}^{2 \pi}\left[1+\frac{g^{\prime}(\xi)}{g\left(\theta_{0}\right)}\left(\theta-\theta_{0}\right)\right]^{n / 2} d \theta\right\}^{1 / n} \\
& \geqslant\|f(z)\|_{\infty}\left\{\frac{1}{2 \pi} \int_{\theta_{0}-1 / n}^{\theta_{0}+1 / n}\left[1+\frac{g^{\prime}(\xi)}{g\left(\theta_{0}\right)}\left(\theta-\theta_{0}\right)\right]_{n / 2} d \theta\right\}_{1 / n} \\
& \geqslant\|f(z)\|_{\infty}\left\{\frac{1}{2 \pi} \int_{\theta_{0}-1 / n}^{\theta_{0}+1 / n}\left[1-\frac{\left\|g^{\prime}(\theta)\right\|_{\infty}}{g\left(\theta_{0}\right)} \cdot \frac{1}{n}\right] d \theta\right\}^{1 / 2} \\
&=\|f(z)\|_{\infty}\left(\frac{1}{\pi n}\right)^{1 / n}\left[1-\frac{\left\|g^{\prime}(\theta)\right\|_{\infty}}{g\left(\theta_{0}\right)} \cdot \frac{1}{n}\right]^{1 / 2} \\
&=\|f(z)\|_{\infty} \cdot \exp \left\{\frac{1}{n} \log \frac{1}{n \pi}\right\} \cdot\left[1+O\left(\frac{1}{n}\right)\right] \\
&=\|f(z)\|_{\infty} \cdot\left[1-\frac{\log \pi n}{n}+O\left(\frac{\log ^{2} n}{n^{2}}\right)\right]\left[1+O\left(\frac{1}{n}\right)\right]
\end{aligned}
$$

This completes the proof of the lemma.
Proof of Theorem 1: By (11) and Lemma 1, we know that

$$
\begin{equation*}
J(n)=\|C(z)\|_{2 n} . \tag{25}
\end{equation*}
$$

By the Hölder inequality, it is easy to see that the $L^{p}$-norm $\|\cdot\|_{p}$ is monotonically increasing in $p$, and hence $J(n)$ is monotonically increasing in $n$. The other results follow from (25) and Lemma 2.

## IV. Example Studies

To illustrate the algorithm works, two examples are studied. They are as follows:

$$
\text { i) } C(z)=1-z-z^{2}
$$

ii) $C(z)=1+2 z+3 z^{2}$.

It is easy to show in example ii) that $\|C(z)\|_{\infty}=6$. However, it is not straightforward to see in example i) that $\|C(z)\|_{\infty}=\sqrt{5}$. After 1500 iterations, the $H_{\infty}$-norm is approximated with relative error under 0.00154 in both cases, which are depicted in Figs. 1 and 2, respectively.

## V. Conclusions and Remarks

a) The proposed algorithm has itself theoretical interests as well as its application importance. Various algorithms for minimization (maximization) of functions exist [6]-[8], but to the authors' knowledge, theoretical algorithms for computing the $H^{\infty}$-norm have not yet been studied elsewhere.
b) It is interesting to note that the principal part of the relative error of the algorithm is independent of the polynomial $C(z)$ (i.e., $(\log n) / 2 n)$. Furthermore, the error is monotonically decreasing to zero. So, for a

given relative error, we can roughly decide the iteration step $n$ to achieve the desired accuracy.
c) In this note, we have only considered the scalar polynomial case. Of course, for a given stable scalar rational function, one can first approximate it by an $r$ th-order polynomial (with exponential decaying error $0\left(\lambda^{r}\right), 0<\lambda<1$ ) and then use the above method to approximate the $H^{\infty}$-norm of the rational function. It is desirable to extend the results of this note to the general matrix transfer function case.

## References

[1] G. Zames and B. A. Francis, "Feedback, minimax sensitivity, and optimal robustness,' IEEE Trans. Automat. Contr., vol. AC-28, pp. 585-601, 1983.
[2] B. A. Francis, "Notes on $H^{\infty}$-optimal linear feedback system," Lecture Notes, Dep. Elec. Eng., Linköping Univ., Linköping, Sweden, 1983.
[3] B. D. O. Anderson and Y. Liu, "Controller reduction: Concepts and approaches," in Proc. Amer. Contr. Conf., MN, 1987, pp. 1-9.
[4] L. Ljung, "On positive real transfer function and the convergence of some recursive schemes," IEEE Trans. Automat. Contr., vol. AC-22, pp. 539-55t, 1977.
[5] H. F. Chen and L. Guo, "Convergence rate of least squares identification and adaptive control for stochastic systems," Int. J. Contr., vol. 44, no. 5, pp. 14591476, 1986.
[6] R. P. Brent, Algorithms for Minimization without Derivatives. Englewood Cliffs, NJ: Prentice-Hall, 1973.
[7] G. P. Szego, Minimization Algorithms: Mathematical Theories and Computer Results. New York: Academic, 1972.
[8] M. R. Osborne, Finite Algorithms in Optimization and Data Analysis. New York: Wiley, 1985.


[^0]:    Manuscript received October 6, 1987; revised January 28, 1988
    The authors are with the Department of Systems Engineering, Research School of Physical Sciences, Australian National University, Canberra, Australia.
    IEEE Log Number 8823316.

